

# Problemen

| Problem Section

**Edition 2020-2** We received solutions from Rik Biel and Thijmen Krebs.

**Problem 2020-2/A** (proposed by Hendrik Lenstra)

Let  $R$  be a ring, and write  $R[[X, X^{-1}]]$  for the set of formal expressions  $\sum_{i \in \mathbb{Z}} a_i X^i$  with all  $a_i \in R$ .

- a. Suppose that  $R[[X, X^{-1}]]$  has a ring structure with the following three properties.
- The sum is given by  $(\sum_{i \in \mathbb{N}} a_i X^i) + (\sum_{i \in \mathbb{Z}} b_i X^i) = \sum_{i \in \mathbb{Z}} (a_i + b_i) X^i$ ,
  - For two formal power series in  $X$ , the product is the regular product of power series, and likewise for two formal power series in  $X^{-1}$ .
  - For  $1 := X^0$ , we have  $X \cdot X^{-1} = 1$ .

Prove that  $R$  is the zero ring.

- b. Prove that for every ring  $R$ , there exists a ring structure on  $R[[X, X^{-1}]]$  satisfying properties I and II.

**Solution** We received a solution from Thijmen Krebs.

For part a, we calculate

$$\begin{aligned} \left( \sum_{i > 0} X^i \right) \cdot \left( \sum_{i < 0} X^i \right) &= \left( \sum_{i \geq 0} X^i \right) \cdot X \cdot X^{-1} \cdot \left( \sum_{i \leq 0} X^i \right) \\ &= \left( \sum_{i \geq 0} X^i \right) \left( \sum_{i \leq 0} X^i \right) \\ &= 1 + \left( \sum_{i < 0} X^i \right) + \left( \sum_{i > 0} X^i \right) + \left( \sum_{i > 0} X^i \right) \cdot \left( \sum_{i < 0} X^i \right) \\ &= \sum_{i \in \mathbb{Z}} X^i + \left( \sum_{i > 0} X^i \right) \cdot \left( \sum_{i < 0} X^i \right). \end{aligned}$$

It follows that the formal power series  $\sum_{i \in \mathbb{Z}} X^i$  equals 0, and in particular,  $1 = 0$ . Hence  $R$  is the zero ring.

For part b, we define

$$\left( \sum_{i \in \mathbb{Z}} a_i X^i \right) \cdot \left( \sum_{i \in \mathbb{Z}} b_i X^i \right) = \left( \sum_{i \geq 0} a_i X^i \right) \left( \sum_{i \geq 0} b_i X^i \right) + \left( \sum_{i \leq 0} a_i X^i \right) \left( \sum_{i \leq 0} b_i X^i \right) - a_0 b_0.$$

This satisfies property (ii), has multiplicative identity 1, and is associative and distributive because multiplication in  $R[[X]]$  and  $R[[X^{-1}]]$  is associative and distributive.

Note that this ring is isomorphic to  $R[[X, Y]] / (XY)$ .

**Problem 2020-2/B** (proposed by Onno Berrevoets)

Let  $n \geq 1$  be an integer, and let  $p_1, \dots, p_{n-1}$  be pairwise distinct prime numbers. Suppose that  $(v_1, \dots, v_n)^\top \in \mathbb{Z}^n$  is a non-trivial element of the kernel of

$$\begin{pmatrix} 1^{p_1} & 2^{p_1} & \dots & n^{p_1} \\ 1^{p_2} & 2^{p_2} & \dots & n^{p_2} \\ \vdots & \vdots & \ddots & \vdots \\ 1^{p_{n-1}} & 2^{p_{n-1}} & \dots & n^{p_{n-1}} \end{pmatrix}.$$

Prove that

$$\max_k |v_k| \geq \frac{2}{n^2 + n} \prod_{i=1}^{n-1} p_i.$$

**Solution** We received no solutions. The proof provided is by Daan and Onno.

Suppose by contradiction that  $\max_k |v_k| < \frac{2}{n^2 + n} \prod_i p_i$ . For every  $i$  the  $i$ -th row vector equals  $(1, 2, \dots, n)$  modulo  $p_i$ , hence  $\sum_k k v_k \equiv 0 \pmod{p_i}$  for every  $i$ . It follows that  $\sum_k k v_k = 0 \pmod{\prod_i p_i}$ , and thus  $\sum_k k v_k = 0$  by the assumed inequality on  $\max_k |v_k|$ . We conclude that  $(1, 2, \dots, n)$  is perpendicular to  $(v_1, \dots, v_n)$ . Now  $(v_1, \dots, v_n)^\top$  is in the kernel of the matrix.

# Oplösungen

| Solutions

$$\begin{pmatrix} 1 & 2 & \dots & n \\ 1^{p_1} & 2^{p_1} & \dots & n^{p_1} \\ 1^{p_2} & 2^{p_2} & \dots & n^{p_2} \\ \vdots & \vdots & \ddots & \vdots \\ 1^{p_{n-1}} & 2^{p_{n-1}} & \dots & n^{p_{n-1}} \end{pmatrix}.$$

Hence, this matrix is singular, and the rows are thus linearly dependent over  $\mathbb{Q}$ . Let  $S = \{1, p_1, p_2, \dots, p_{n-1}\}$  and let  $(\lambda_s)_{s \in S} \in \mathbb{Q}^n$  be such that

$$\sum_{s \in S} \lambda_s (1^s, 2^s, \dots, n^s) = (0, 0, \dots, 0).$$

Then the non-zero polynomial  $\sum_s \lambda_s X^s$  has roots  $1, 2, \dots, n$  and at most  $n$  non-zero coefficients.

We claim that such a polynomial does not exist. More precisely, we claim that any non-zero polynomial in  $\mathbb{R}[X]$  with at least  $n$  distinct positive roots has at least  $n + 1$  non-zero coefficients. We do this by induction.

The case  $n = 0$  is trivial. For  $N > 0$ , assume by induction that any non-zero polynomial with at least  $N - 1$  distinct positive roots has at least  $N$  non-zero coefficients. Suppose that  $f$  is a non-zero polynomial with at least  $N$  positive roots. Without loss of generality, assume the constant term of  $f$  is non-zero (dividing by  $X$  does not change the number of terms or the number of positive roots of  $f$ ), and moreover,  $f$  is non-constant. Suppose  $f$  has positive roots  $a_1 < a_2 < \dots < a_N$ . Then the derivative of  $f$  has a positive root between  $a_i$  and  $a_{i+1}$  for all  $i \in \{1, 2, \dots, N - 1\}$ , so it has at least  $N - 1$  distinct positive roots. By the induction hypothesis, it has at least  $N$  non-zero coefficients. Since  $f$  had a non-zero constant term by assumption, it follows that  $f$  must have had at least  $N + 1$  non-zero coefficients to begin with.

This proves the claim, and shows in particular that there is no non-zero polynomial with at most  $n$  non-zero terms with roots  $1, 2, \dots, n$ . This contradiction shows that we cannot have  $\max_k |v_k| < \frac{2}{n^2+n} * \prod_i p_i$ , as was to be shown.

**Problem 2020-2/C** (proposed by Onno Berrevoets)

Let  $n, m, k \geq 2$  be positive integers.  $n$  students will attend a multiple-choice exam containing  $mk$  questions and each questions has  $k$  possible answers. A student passes the exam precisely when he/she answers at least  $m + 1$  questions correctly.

- Suppose that  $n = 2k$ . Show that the students can coordinate their answers such that it is guaranteed that at least one student passes the exam.
- Suppose that  $n = 2k - 1$ . Does there exist a  $k$  for which the students can coordinate their answers such that it is guaranteed that at least one student passes the exam?

**Solution** We received solutions from Rik Biel and Thijmen Krebs. This solution is based on the one by Thijmen.

We describe a strategy for the students by giving  $mk$  ordered partitions of  $\{1, \dots, n\}$  into  $k$  parts, indicating that student  $s$  picks answer  $a$  on question  $q$  if and only if  $s$  is in the  $a$ -th part of the  $q$ -th partition.

For part a, we take

$$\begin{aligned} &\{1, 2\}, \{3, 4\}, \dots, \{n - 1, n\} \quad (\text{first } mk - 1 \text{ questions}), \\ &\{2, 3\}, \{4, 5\}, \dots, \{n, 1\} \quad (\text{final question}). \end{aligned}$$

Note that student  $2i$  is in the  $i$ -th part of each partition, and that student  $2i - 1$  is in the  $i$ -th part of each partition except the last, in which case she is in the  $i - 1$ -th part (taken cyclically). In other words, student  $2i$  answers  $i$  on each question, while student  $2i - 1$  answers  $i$  on all but the last question, on which she answers  $i - 1$ .

The only way in which none of the students passes on the first  $mk - 1$  questions is if all but one of the student pairs  $\{2i - 1, 2i\}$  get exactly  $m$  correct answers and the final pair gets exactly  $m - 1$  correct answers on these questions. Assuming this is the case, then if the correct answer to the final question is  $i$ , we find that both students  $2i$  and  $2i + 1$  get

another correct answer, and at least one of these two already had  $m$  correct answers on the first  $mk - 1$  questions. Thus, at least one student passes.

For part b, there exists such  $k$  if  $k > m + 2$ . Take

$$\begin{aligned} \{1, 2\}, \{3, 4\}, \dots, \{n-2, n-1\}, \{n\} & \quad m(k-1) - 1 \text{ times,} \\ \{2, 3\}, \{4, 5\}, \dots, \{n-1, 1\}, \{n\} & \quad m + 1 \text{ times.} \end{aligned}$$

Without loss of generality, we assume student  $n$  is correct on the first  $i$  questions and the last  $m-i$  questions with  $0 \leq i \leq m$ ; if she is correct more often, we are done, and otherwise, the situation can only improve for the remaining students. Then there are  $m(k-1) - (i+1)$  questions of the first kind and  $1 \leq i+1 < k-1$  questions of the second kind for the other students to score points on. On each of these questions, precisely two students score points, for a total of  $2m(k-1)$  points distributed over  $2(k-1)$  students. This means that no student passes if and only if each student scores exactly  $m$  points.

The only way to avoid passing a student on the remaining  $m(k-1) - (i+1)$  questions of the first kind is for no fewer than  $k-1 - (i+1)$  of the  $k-1$  pairs  $\{1, 2\}, \dots, \{n-2, n-1\}$  to have exactly  $m$  correct answers; otherwise, the total number of questions of this kind can at most be  $(k-1 - (i+2))m + (i+2)(m-1) = m(k-1) - (i+2)$ . The remaining  $i+1$  pairs have in total  $2((i+1)m - (i+1))$  correct answers. In particular, since  $1 \leq i+1 < k-1$ , there is at least one pair of students with exactly  $m$  points and at least one pair with strictly less than  $m$  points scored on the questions of the first kind. Working cyclically, we can therefore assume that the pair  $\{2j-1, 2j\}$  has exactly  $m$  correct answers and  $\{2j+1, 2j+2\}$  has strictly less than  $m$  correct answers in questions of the first kind. We find that student  $2j$  cannot have any correct answer in the final questions, and since she answers in the same way as student  $2j+1$ , neither can student  $2j+1$ . But then student  $2j+1$  will in total score strictly less than  $m$  points, meaning that at least one other student will score strictly more than  $m+1$  points. In other words, this strategy guarantees that at least one student passes.